## **CBMC**

#### A Bounded Model Checker for C programs

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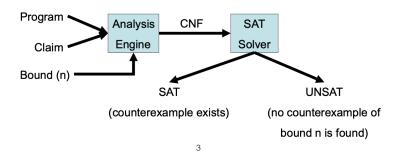
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(First part is based on slides by Arie Gurfinkel.)

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## Bug catching with CBMC

- Main Ideia: Given a program and a claim use a SAT solver to find whether there exists an execution that violates the claim.
  - Arbitrary ANSI-C programs (with bitvector arithmetic, dynamic memory, pointers, type casts, etc).
  - Simple Safety Claims (such as, array bound checks, pointer checks, division by zero, arithmetic overflow, user supplied assertions).
- Loops must be bounded (i.e., the analysis is incomplete).
- The analysis is completely automated.



#### **CBMC** — C Bounded Model Checker

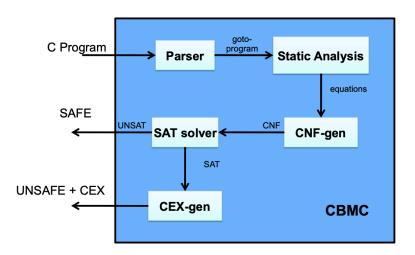
- Developed at CMU by Daniel Kroening (Oxford) and Edmund Clarke (CMU).
- Available at: http://www.cprover.org/cbmc/
- CBMC demonstrates the violation of assertions in C programs, or proves safety of the assertions under a given bound.
- CBMC implements a bit-precise translation of an input C program, annotated with assertions and with loops unrolled to a given depth, into a formula. If the formula is satisfiable, then an execution leading to a violated assertion exists.
- CBMC is not able to prove correctness for programs with unbounded loops in general, but is very useful for bug catching.

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### How does CBMC work

- Transform a program into a set of equations.
  - 1. Simplify control flow
  - 2. Unwind all the loops
  - 3. Convert into Static Single Assignment (SSA) form
  - 4. Convert into equations
  - 5. Bit-blast
- Solve the equations with a SAT solver.
- Convert the SAT assignment (if any) into a counterexample.

### How does CBMC work



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## Loop unwinding

- All loops are unwound.
  - different unwinding bounds for different loops can be used
  - to check whether unwinding is sufficient special "unwinding assertion" claims are added
- If a program satisfies all of its claims and all unwinding assertions then it is correct!
- Same for backward goto jumps and recursive functions.

## Control flow simplifications

· All side effect are removed

```
e.g., j=i++ becomes j=i;i=i+1
```

· Control flow is made explicit

```
continue, break replaced by goto
```

· All loops are simplified into one form

```
for, do-while replaced by while
```

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### Loop unwinding

```
void f(...) {
                                       void f(...) {
                                                            void f(...) {
void f(...) {
                                                              if(cond) {
                                         if(cond) {
                      if(cond) {
  while(cond) {
                                           Body;
                                                                Body;
    Body;
                                                                if(cond) {
                                           if(cond) {
                        while(cond) {
                                                                  Body;
                          Body;
  Remainder;
                                                                  if(cond) {
                                             while(cond) {
                                                                    Body:
                                                                    assert(!cond):
                      Remainder;
                                                                               Unwinding
                                         Remainder;
                                                                               assertion
                                                               Remainder
```

This allows us to prove that we have done enough unwinding.

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### Example of a sufficient loop unwinding

```
void f(...) {
    j = 1
    while (j <= 2)
    j = j + 1;
    Remainder;
}</pre>
```

unwind = 3

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#### Transforming loop-free programs into equations

When a variable is assigned multiple times, use a new variable for the LHS of each assignment.

This is called Static Single Assignment (SSA) form.

#### Program



#### SSA Program

```
x_1 = x_0 + y_0;

x_2 = x_1 * 2;

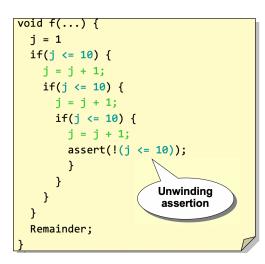
a_1[i_0] = 100;
```

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### Example of a insufficient loop unwinding

```
void f(...) {
   j = 1
   while (j <= 10)
    j = j + 1;
   Remainder;
}</pre>
```

unwind = 3



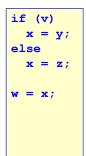
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### Transforming loop-free programs into equations

What about conditionals?

For each join point, add new variables with selectors.

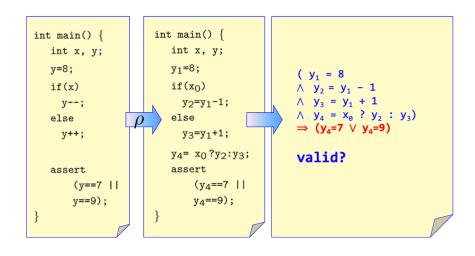
#### Program



#### SSA Program

```
if (v<sub>0</sub>)
  x<sub>1</sub> = y<sub>0</sub>;
else
  x<sub>2</sub> = z<sub>0</sub>;
  x<sub>3</sub> = v<sub>0</sub> ? x<sub>1</sub> : x<sub>2</sub>;
  w<sub>1</sub> = x<sub>3</sub>;
```

## Example

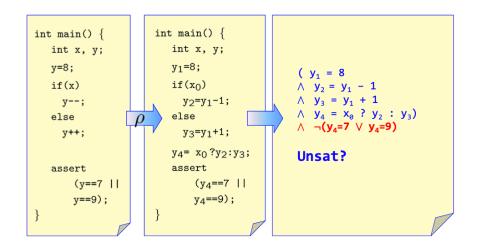


## Bit-basting

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- To check satisfiability of the formulas the usual approach is to replace the arithmetic operators by circuit equivalents to obtain a propositional formula, which is then passed to a SAT solver.
- This approach is called 'bit-blasting' or 'bit-flattening', as the word-level structure is lost.
- CBMC also supports SMT solvers as back-ends. In this case a bit-vector theory is used.

## Example



Extend C programming language with 3 modelling features.

#### **Assertions**

assert (e) - aborts an execution when e is false, no-op otherwise.

#### Non-determinism

nondet\_int() - returns a non-deterministic integer value.

#### **Assumptions**

**CPROVER** assume (e) - "ignores" execution when e is false.

Restricts the program traces that are considered and allows assume-guarantee reasoning.

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## From programming to modelling

## Assume-Guarantee reasoning

(A) Is foo correct assuming p is not NULL?

```
int foo (int* p) { __CPROVER_assume(p!=NULL); ... }
```

(G)Is foo guaranteed to be called with a non-NULL argument?

```
void main(void) {
    ...
    assert (x!=NULL);// foo(x);
    ...
    assert (y!=NULL); //foo(y);
    ...}
```

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# **CBMC** by example

### Dangers of unrestricted assumptions

Assumptions can lead to vacuous satisfaction.

This program is passed by CBMC!

```
if (x > 0) {
   __CPROVER_assume (x < 0);
   assert (0); }</pre>
```

Assumptions can be used to restrict the program traces that are considered.

```
x = nondet_int ();
y = nondet_int ();
__CPROVER_assume (x < y);</pre>
```

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## Assertions [ex1.c]

CBMC checks assertions as defined by the ANSI-C standard. The assert statement takes a Boolean condition, and CBMC checks that this condition is true for all runs of the program.

```
void main (void)
{
  int x;
  int y=8, z=0, w=0;

  if (x)
    z = y - 1;
  else
    w = y + 1;

  assert (z == 7 || w == 9);
}
```

```
$ cbmc ex1.c
$ cbmc ex1.c --show-vcc
```

### ex1.c outcome

#### \$ cbmc ex1.c

```
CBMC version 5.10 (cbmc-5.10) 64-bit x86_64 macos
Parsing ex1.c
Converting
Type-checking ex1
file ex1.c line 11 function main: function `assert' is not declared
Generating GOTO Program
Adding CPROVER library (x86_64)
Removal of function pointers and virtual functions
Generic Property Instrumentation
Running with 8 object bits, 56 offset bits (default)
Starting Bounded Model Checking
size of program expression: 43 steps
simple slicing removed 2 assignments
Generated 1 VCC(s), 1 remaining after simplification
Passing problem to propositional reduction
converting SSA
Running propositional reduction
Post-processing
Solving with MiniSAT 2.2.1 with simplifier
141 variables, 39 clauses
SAT checker inconsistent: instance is UNSATISFIABLE
Runtime decision procedure: 0.00294518s
[main.assertion.1] assertion z == 7 \parallel w == 9: SUCCESS
** 0 of 1 failed (1 iteration)
VERIFICATION SUCCESSFUL
```

### Alternatively: use SMT solver

#### \$ cbmc -z3 ex1.c

```
(...)
Generated 1 VCC(s), 1 remaining after simplification
Passing problem to SMT2 QF_AUFBV using Z3
converting SSA
Running SMT2 QF_AUFBV using Z3
Runtime decision procedure: 0.0279709s

** Results:
[main.assertion.1] assertion z == 7 || w == 9: SUCCESS

** 0 of 1 failed (1 iteration)
VERIFICATION SUCCESSFUL
```

#### S cbmc -cvc4 ex1.c

```
(...)
Generated 1 VCC(s), 1 remaining after simplification
Passing problem to SMT2 QF_AUFBV using CVC4
converting SSA
Running SMT2 QF_AUFBV using CVC4
Runtime decision procedure: 0.0147195s

** Results:
[main.assertion.1] assertion z == 7 || w == 9: SUCCESS

** 0 of 1 failed (1 iteration)
VERIFICATION SUCCESSFUL
```

#### ex1.c outcome

#### \$ cbmc ex1.c --show-vcc

#### ex2.c

```
void main (void)
{
  int x;
  int y=8, z=0, w=0;

  if (x)
    z = y - 1;
  else
    w = y + 1;

  assert (z == 5 || w == 9);
}
```

```
$ cbmc ex2.c
$ cbmc ex2.c --trace
```

### ex2.c outcome

#### \$ cbmc ex2.c

```
(...)
Solving with MiniSAT 2.2.1 with simplifier
141 variables, 50 clauses
SAT checker: instance is SATISFIABLE
Runtime decision procedure: 0.000694225s

** Results:
[main.assertion.1] assertion z == 5 || w == 9: FAILURE

** 1 of 1 failed (1 iteration)
VERIFICATION FAILED
```

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#### ex2.c outcome

#### \$ cbmc ex2.c --trace

### ex2.c outcome

ex3.c

CBMC can ignore user assertions.

```
void main (void)
{
  int x, y;

  x = x + y;
  if (x != 3) x = 2;
  else x++;

  assert (x <= 3);
}</pre>
```

```
$ cbmc ex3.c --show-vcc
$ cbmc ex3.c
$ cbmc ex3.c --no-assertions
$ cbmc ex3.c --no-assertions --show-vcc
```

## Checking overflow

But the code can be automatically instrumented.

```
$ cbmc ex3.c --signed-overflow-check --no-assertions --trace
```

```
State 17 file ex3.c line 3 function main thread 0

x=-1610612735 (10100000 00000000 00000000 00000001)

State 18 file ex3.c line 3 function main thread 0

y=-2147483648 (10000000 00000000 00000000 00000000)

Violated property:
file ex3.c line 5 function main
arithmetic overflow on signed + in x + y
!overflow("+", signed int, x, y)

** 1 of 2 failed (2 iterations)
VERIFICATION FAILED
```

## Seeing the properties

\$ cbmc ex3.c --signed-overflow-check --show-properties

```
Property main.overflow.1:
    file ex3.c line 5 function main
    arithmetic overflow on signed + in x + y
    !overflow("+", signed int, x, y)

Property main.overflow.2:
    file ex3.c line 7 function main
    arithmetic overflow on signed + in x + 1
    !overflow("+", signed int, x, 1)

Property main.assertion.1:
    file ex3.c line 9 function main
    assertion x <= 3
    x <= 3
```

#### Workflow

• Internally CBMC runs goto-cc to produce a representation of the control flow graph of the program.

```
$ goto-cc ex3.c -o ex3.gb
```

 Then the instrumentation tool goto-instrument automatically add assertions to be checked.

```
$ goto-instrument --signed-overflow-check ex3.gb
ex3.instr.qb
```

· And finally the assertions are checked.

```
$ cbmc ex3.instr.gb
```

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### Seeing the instrumented code

```
$ cbmc ex3.c --signed-overflow-check --show-goto-functions
```

```
main /* main */
        // 0 file ex3.c line 3 function main
       signed int x:
       // 1 file ex3.c line 3 function main
        signed int y;
        // 2 file ex3.c line 5 function main
        ASSERT !overflow("+", signed int, x, y) // arithmetic overflow on signed + in x + y
       // 3 file ex3.c line 5 function main
        // 4 file ex3.c line 6 function main
       IF !(x != 3) THEN GOTO 1
       // 5 file ex3.c line 6 function main
       // 6 file ex3.c line 6 function main
       GOTO 2
       // 7 file ex3.c line 7 function main
     1: ASSERT !overflow("+", signed int, x, 1) // arithmetic overflow on signed + in x + 1
        // 8 file ex3.c line 7 function main
        // 9 file ex3.c line 9 function main
     2: ASSERT x \le 3 // assertion x \le 3
        // 10 file ex3.c line 10 function main
        // 11 file ex3.c line 10 function main
        // 12 file ex3.c line 10 function main
        END_FUNCTION
```

## Entrypoints [ex4.c]

```
int fun (int a, int b)
{
  int c = a+b;

  if (a>0 || b>0)
    c = 1/(a+b);
  return c;
}
```

```
$ cbmc ex4.c
$ cbmc ex4.c --function fun
$ cbmc ex4.c --function fun --div-by-zero-check
$ cbmc ex4.c --function fun --div-by-zero-check --trace
```

#### ex5.c

```
void main ()
{
   char c;
   long l;
   int i;

   l = c = i;
   assert (l==i);
}
```

```
$ cbmc ex5.c
$ cbmc ex5.c --trace
```

## Checking division by zero

\$ cbmc ex4.c --function fun --div-by-zero-check --trace [fun.division-by-zero.1] division by zero in 1 / (a + b): FAILURE Trace for fun.division-by-zero.1: State 23 file ex4.c line 1 thread 0 a=-1073741808 (11000000 00000000 00000000 00010000) State 24 file ex4.c line 1 thread 0 b=1073741808 (00111111 11111111 11111111 11110000) State 25 file ex4.c line 3 function fun thread 0 C=0 (00000000 00000000 00000000 00000000) State 26 file ex4.c line 3 function fun thread 0 c=0 (00000000 00000000 00000000 00000000) Violated property: file ex4.c line 6 function fun division by zero in 1 / (a + b)!(a + b == 0)\*\* 1 of 1 failed (1 iteration) VERIFICATION FAILED

#### ex5.c outcome

\$ cbmc ex5.c --trace

```
** Results:
[main.assertion.1] assertion l == (signed long int)i: FAILURE
Trace for main.assertion.1:
State 17 file ex5.c line 3 function main thread 0
 c=0 (00000000)
State 18 file ex5.c line 4 function main thread 0
 State 19 file ex5.c line 5 function main thread 0
 i=262144 (00000000 00000100 00000000 00000000)
State 20 file ex5.c line 7 function main thread 0
 c=0 (00000000)
State 21 file ex5.c line 7 function main thread 0
 Violated property:
 file ex5.c line 8 function main
 assertion l == (signed long int)i
 l == (signed long int)i
```

## Array bounds [ex6.c]

```
int puts (const char *s);
int main (int argc, char **argv)
{
  int i;
  if (argc >= 1)
    puts (argv[2]);
}
```

```
$ cbmc ex6.c
$ cbmc ex6.c --bounds-check --pointer-check
$ cbmc ex6.c --bounds-check --pointer-check --trace
```

## Array bounds [ex7.c]

```
int puts (const char *s);
int main (int argc, char **argv)
{
  int i;

  if (argc >= 2)
    puts (argv[2]);
}
```

\$ cbmc ex7.c --bounds-check --pointer-check

```
(...)
[main.pointer_dereference.7] dereference failure: invalid integer address
in argv[(signed long int)2]: SUCCESS

** 0 of 7 failed (1 iteration)
VERIFICATION SUCCESSFUL
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```

### ex6.c outcome

\$ cbmc ex6.c --bounds-check --pointer-check --trace

```
Trace for main.pointer_dereference.6:
State 17 thread 0
 INPUT arac: 1 (00000000 00000000 00000000 00000001)
State 18 thread 0
 State 21 file ex6.c line 3 thread 0
 argc=1 (00000000 00000000 00000000 00000001)
State 22 file ex6.c line 3 thread 0
 State 23 file ex6.c line 5 function main thread 0
 i=0 (00000000 00000000 00000000 00000000)
Violated property:
 file ex6.c line 8 function main
 dereference failure: pointer outside object bounds in argv[(signed long int)2]
 16l + POINTER_OFFSET(argv) >= 0l && OBJECT_SIZE(argv) >= 24ul + (unsigned long
int)POINTER_OFFSET(argv) || DYNAMIC_OBJECT(argv)
** 1 of 7 failed (2 iterations)
VERIFICATION FAILED
```

### ex8.c

```
int array[10];
int sum ()
{
  unsigned i, sum;

  sum = 0;
  for (i = 0; i <= 10; i++)
     sum += array [i];
}</pre>
```

```
$ cbmc ex8.c --function sum
$ cbmc ex8.c --function sum --bounds-check
$ cbmc ex8.c --function sum --bounds-check --trace
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```

#### ex8.c outcome

\$ cbmc ex8.c --function sum --bounds-check --trace

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#### ex9.c outcome

```
$ cbmc ex9.c --function binsearch
--bounds-check --pointer-check
```

**CBMC** does not stop! The loop is being infinitely unwound. We must provide the number of iterations to be unwound.

```
$ cbmc ex9.c --function binsearch
--bounds-check --pointer-check
--unwind 4
```

```
(...)
[binsearch.array_bounds.4] array `a' upper bound in a[(signed long int)middle]: SUCCESS
** 0 of 4 failed (1 iteration)
VERIFICATION SUCCESSFUL
```

The above verification simply means that no array bounds are violated in the first 4 iterations os the loop!

## Loop unwinding [ex9.c]

```
int binsearch (int x)
{
  int a[16];
  signed low = 0, high = 16;

while (low < high) {
    signed middle = low + ((high - low) >> 1);
    if (a[middle] < x) high = middle;
    else if (a[middle] > x) low = middle + 1;
    else return middle;
  }
  return -1;
}
```

```
$ cbmc ex9.c --function binsearch
--bounds-check --pointer-check
```

### Unwinding assertion

To see if the entire set of possible executions is being covered, we must generate *unwinding assertions*.

```
(...)
[binsearch.array_bounds.4] array `a' upper bound in a[(signed long int)middle]: SUCCESS
[binsearch.unwind.0] unwinding assertion loop 0: FAILURE

** 1 of 5 failed (2 iterations)
VERIFICATION FAILED
```

## Unwinding assertion

The failure of the "unwinding assertion" means that it is not guaranteed that the number k of iterations given as parameter will be sufficient, i.e. some execution path may run through n>k iterations.

In this case it suffices to increase k.

```
$ cbmc ex9.c --function binsearch
--bounds-check --pointer-check
--unwind 6 --unwinding-assertions
```

```
(...)
** Results:
[binsearch.array_bounds.1] array `a' lower bound in a[(signed long int)middle]: SUCCESS
[binsearch.array_bounds.2] array `a' upper bound in a[(signed long int)middle]: SUCCESS
[binsearch.array_bounds.3] array `a' lower bound in a[(signed long int)middle]: SUCCESS
[binsearch.array_bounds.4] array `a' upper bound in a[(signed long int)middle]: SUCCESS
[binsearch.unwind.0] unwinding assertion loop 0: SUCCESS

** 0 of 5 failed (1 iteration)
VERIFICATION SUCCESSFUL
```

## Unbounded loops [ex11.c]

CBMC can also be used for programs with unbounded loops.

```
int sumqq (int x)
{
    short int i, s;

    s = 0;
    for (i = 0; i <= x; i++)
        s += i+i;
    return s;
}</pre>
```

```
$ cbmc ex11.c --function sumqq --signed-overflow-check --unwind 100
```

## Bounded loops [ex10.c]

CBMC checks if enough unwinding is done.

```
int sumq (void)
{
    short int i, s;

s = 0;
    for (i = 0; i <= 10; i++)
        s *= i*i;
    return s;
}</pre>
```

\$ cbmc ex10.c --function sumq --signed-overflow-check

```
[sumq.overflow.3] line 6 arithmetic overflow on signed + in i + 1: SUCCESS
[sumq.overflow.1] line 7 arithmetic overflow on signed * in (signed int)i * (signed int)i: SUCCESS
[sumq.overflow.2] line 7 arithmetic overflow on signed * in s * (signed short int)((signed int)i * (signed int)i): SUCCESS

** 0 of 3 failed (1 iteration)
VERIFICATION SUCCESSFUL

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```

## Unbounded loops [ex11.c]

```
$ cbmc ex11.c --function sumqq --signed-overflow-check --unwind 100
```

```
** Results:
(...)
[sumqq.overflow.3] arithmetic overflow on signed + in i + 1: SUCCESS

** 0 of 3 failed (1 iteration)
VERIFICATION SUCCESSFUL
```

In this case CBMC is used for bug hunting only. CBMC does not attempt to find all bugs. If you increase the bound you can find a bug.

```
$ cbmc ex11.c --function sumqq --signed-overflow-check --unwind 200
```

```
[sumqq.overflow.3] line 6 arithmetic overflow on signed + in i + 1: SUCCESS
[sumqq.overflow.1] line 7 arithmetic overflow on signed + in (signed int)i + (signed int)i: SUCCESS
[sumqq.overflow.2] line 7 arithmetic overflow on signed + in s + (signed short int)((signed int)i + (signed int)i): FAILURE

** 1 of 3 failed (2 iterations)
VERIFICATION FAILED

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```

## Inlining [ex12.c]

CBMC supports functions by inlining, and preserves the locality of the parameters and the non-static local variables by renaming.

```
int f() {
    static int s=0;
    s++;
    return s;
}
int g() {
    int l=0;
    l++;
    return l;
}
```

```
int h() {
   int x=10;
   x += f() + g();
   return x;
}

void main(void)
{
   assert(f()==1); // first call to f
   assert(g()==1); // second call to f
   assert(g()==1); // second call to g
   assert(g()==14);
}
```

#### S cbmc ex12.c

\*\* 0 of 1 failed (1 iteration)
VERIFICATION SUCCESSFUL
Generated 5 VCC(s), 0 remaining after simplification
VERIFICATION SUCCESSFUL

## Recursion [ex13.c]

\$ cbmc ex13.c --function fact --signed-overflow-check
--unwind 5

```
[fact.overflow.1] line 5 arithmetic overflow on signed - in n - 1: FAILURE
[fact.overflow.2] line 5 arithmetic overflow on signed * in n * return_value_fact: SUCCESS
** 1 of 2 failed (2 iterations)

VERIFICATION FAILED
```

If called from main fact will be inlined and unwound.

There is no need to provide --unwind k switch.

\$ cbmc ex13.c --signed-overflow-check

```
[fact.overflow.1] line 5 arithmetic overflow on signed - in n - 1: SUCCESS
[fact.overflow.2] line 5 arithmetic overflow on signed * in n * return_value_fact: FAILURE

** 1 of 2 failed (2 iterations)
VERIFICATION FAILED
```

## Recursion [ex13.c]

Recursion is implemented by finite unwinding, as done for while loops.

```
int fact(int n)
{
   if (n==0)
     return 1;
   else return n*fact(n-1);
}

void main(void)
{
   fact(20);
}
```

#### \$ cbmc ex13.c --function fact -signed-overflow-check

```
Unwinding recursion fact iteration 1
Unwinding recursion fact iteration 2
Unwinding recursion fact iteration 3
Unwinding recursion fact iteration 4
Unwinding recursion fact iteration .....
```

**CBMC** does not stop! The loop is being infinitely unwound.

We must provide the number of iterations to be unwound.

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## Low level properties [ex14.c]

Nondeterminism can be introduced explicitly into the program by means of functions that begin with the prefix **nondet**\_

```
int nondet_int();
int *p;
int global;

void f (void)
{
  int local = 10;
  int input = nondet_int();

  p = input ? &local : &global;
}
```

```
int main (void)
{
  int z;

  global = 10;
  f ();
  z = *p;
  assert (z==10);
}
```

### ex15.c

```
int fun (int n)
{
  int *p, i, s=0;

  p = malloc(sizeof(int)*n);
  for (i=0; i<++n; i++)
    p[i] = 10*i;
  for (i=0; i<n; i++)
    s += p[i];
  return s;
}</pre>
```

```
int main(void)
{
  printf("%d",fun(8) + 100);
  return 0;
}
```

\$ cbmc ex15.c --bounds-check --pointer-check

The loop is being infinitely unwound.

Why?

VERIFICATION SUCCESSFUL W

Why?

\$ cbmc ex15.c --bounds-check -pointer-check -unwind 10

**VERIFICATION FAILED** 

Why?

### ex17.c

```
void f (unsigned int n)
{
  int *p;
  p = malloc(sizeof(int)*n);
  p[n-1] = 0;
  free(p);
}
```

\$ cbmc ex17.c --function f VERIFICATION SUCCESSFUL

VERIFICATION FAILED Why?

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### ex16.c

```
char *p = "abc";
void fun(unsigned int i)
{
  char ch;
  ch = p[i];
}
```

#### \$ cbmc ex16.c --bounds-check --pointer-check --function fun

```
(...)
[fun.pointer_dereference.6] dereference failure: pointer outside object bounds in p[(signed long int)i]: FAILURE
[fun.pointer_dereference.7] dereference failure: invalid integer address in p[(signed long int)i]: SUCCESS

** 1 of 7 failed (2 iterations)
VERIFICATION FAILED
```

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#### ex18.c

```
void f (int i)
{
  int *p, y;

  p = malloc(sizeof(int)*10);
  if (i) p = &y;
  free(p);
}
```

\$ cbmc ex18.c --function f

VERIFICATION FAILED Why?

### Assume-guarantee reasoning

- In addition to the assert statement, CBMC provides the <u>\_\_CPROVER\_assume</u> statement.
- As an assertion, CPROVER assume takes a Boolean expression.
- The \_\_CPROVER\_assume statement restricts the program traces that are considered and allows assume-guarantee reasoning.
- Intuitively, one can consider the \_\_CPROVER\_assume statement to abort the
  program successfully if the condition is false. If the condition is true, the
  execution continues.

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### ex20.c

```
int nondet_int();
int x, y;

void main (void)
{
    x = nondet_int();
    __CPROVER_assume (x<10);
    y = x+1;
    assert (y>x);
}
```

\$ cbmc ex20.c

VERIFICATION SUCCESSFUL

\$ cbmc ex20.c --show-vcc

```
VERIFICATION CONDITIONS:
(...)

assertion y > x
(...)
{-12} x#1 = 0
{-13} y#1 = 0
{-14} x#2 = nondet_symbol identifier="symex::nondet0"
{-15} -(x#2 ≥ 10)
{-16} y#2 = 1 + x#2
{1} -(x#2 ≥ y#2)
```

### ex19.c

```
int nondet_int();
int x, y;

void main (void)
{
    x = nondet_int();
    y = x+1;
    assert (y>x);
}
```

```
$ cbmc ex19.c

VERIFICATION FAILED Why?

$ cbmc ex19.c --show-vcc
```

## Lab session