Verification of Timed Systems

Renato Neves





Universidade do Minho

Given a system ${\it S}$ and a property φ show that



The choice of which logical language to use for writing φ depends on the underlying computational paradigm

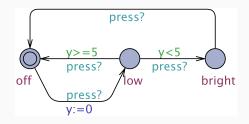
Variant of Computation Tree Logic with two types of formulae

description of state and path properties

Grammar

 Ψ ::= $\ell \mid c \mid$ deadlock \mid not $\Psi \mid \Psi$ or $\Psi \mid \Psi$ and Ψ

We can thus express current locations ℓ , clock constraints $c \in C(C)$, and the presence of deadlocks



Exercise

Write formulae for the following statements

- 1. The lamp is low
- 2. Not off and y > 25
- 3. If it is low or bright then $y \leq 3600$

Grammar

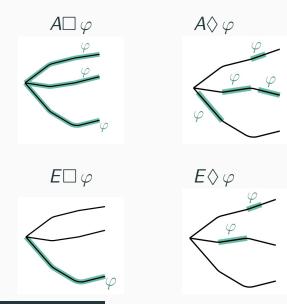
$$\Pi ::= \mathbf{A} \Box \Psi \mid \mathbf{A} \Diamond \Psi \mid \mathbf{E} \Box \Psi \mid \mathbf{E} \Diamond \Psi \mid \Phi \rightsquigarrow \Psi$$

where

- A, E quantify (universally and existentially, resp.) over paths
- quantify (universally and existentially, resp.) over states in a path

Paths can be seen as possible executions

Path Formulae pt. II



$A\Box \varphi$ and $E\Box \varphi$

Something bad will never happen

Examples:

- A nuclear reactor's temperature will never exceed a certain threshold
- We will never reach deadlock
- There is at least one execution in which we never reach deadlock

$E\Diamond \varphi$

Something good can happen

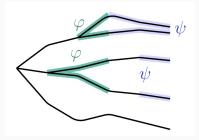
Examples:

- All adventurers reach the other side.
- All adventurers reach the other side in \leq 17 minutes.

•

For all paths if φ holds at some point then ψ will also hold later on

$$\varphi \rightsquigarrow \psi$$



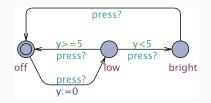
$A \diamondsuit \phi$ and $\phi \rightsquigarrow \psi$ Something good will eventually happen or if something happens then something good will eventually happen Examples:

- Always when pressing the on button the television will eventually turn on
- If the philosopher requests a fork she will eventually get it
- If the plane asks to land it will eventually land

Write the sentences below in CTL

- 1. The system never enters in deadlock
- 2. The location ℓ is reachable
- 3. In all executions we reach location ℓ
- 4. If we reach location ℓ we will inevitably reach location s
- 5. There exists at least one execution where variable i is always below or equal 10
- 6. The two philosophers never eat at the same time

Back to the lamp



Exercise

- 1. The lamp can become bright;
- 2. The lamp will eventually become bright;
- 3. The lamp can never be on for more than 3600s;
- 4. It is possible to never turn on the lamp;
- 5. Whenever the light is bright, the clock y is non-zero;
- 6. Whenever the light is bright it will eventually become off.